

Simulability of open quantum system dynamics

“A dissipative quantum Church-Turing theorem”

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Simulability

What is (efficiently) simulatable?



[1] A. M. Turing, Proc. London Math. Soc. 42 (1937) no. 230, 230–265

Simulability

What is (efficiently) simulatable?



Here: [Quantum many body dynamics](#)

- On a quantum computer?
- On a classical computer?

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 - Dissipative Church-Turing theorem

Preliminaries

Unitary vs. Liouvillian dynamics

Unitary:

equation of motion:
$$\frac{d}{dt}\rho(t) = -i[\mathbf{H}, \rho(t)]$$

Unitary vs. Liouvillian dynamics

Unitary:

Liouvillian:

equation of motion:

$$\frac{d}{dt}\rho(t) = -i[\mathbf{H}, \rho(t)]$$

$$\frac{d}{dt}\rho(t) = \mathcal{L}(\rho(t))$$

Unitary vs. Liouvillian dynamics

Unitary:

Liouvillian:

equation of motion: $\frac{d}{dt}\rho(t) = -i[\mathbf{H}, \rho(t)]$

$$\frac{d}{dt}\rho(t) = \mathcal{L}(\rho(t))$$

time independent: $\rho(t) = e^{-i\mathbf{H}t}\rho(0)e^{i\mathbf{H}t}$

$$\rho(t) = e^{\mathcal{L}t}\rho(0)$$

Unitary vs. Liouvillian dynamics

	Unitary:	Liouvillian:
equation of motion:	$\frac{d}{dt}\rho(t) = -i[\mathbf{H}, \rho(t)]$	$\frac{d}{dt}\rho(t) = \mathcal{L}(\rho(t))$
time independent:	$\rho(t) = e^{-i\mathbf{H}t}\rho(0)e^{i\mathbf{H}t}$	$\rho(t) = e^{\mathcal{L}t}\rho(0)$
time dependent:	"time ordered product integrals"	

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time dependent:	"time ordered product integrals"	

Propagator for $t \geq s \geq 0$

$$\rho(t) = T_{\mathcal{L}}(t, s)(\rho(s))$$

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Liouvillian:

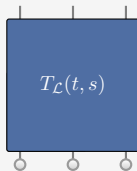
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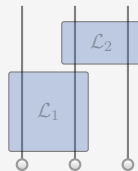
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Distinguishability of propagators

Distinguishability of density matrices:

$$\|\rho - \sigma\|_1 = \max_{0 \leq A \leq \mathbf{1}} \text{tr}(A(\rho - \sigma))$$

Worst case estimate for propagators:

$$\|T - T'\|_{1 \rightarrow 1} := \sup_{\|\rho\|_1=1} \|T(\rho) - T'(\rho)\|_1$$

Trotterization of Liouvillian dynamics

k -local Liouvillian dynamics

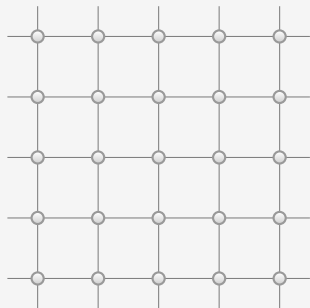
$$\mathcal{L} = \sum_{\Lambda}^K \mathcal{L}_{\Lambda} \quad \mathcal{L}_{\Lambda}(\rho) = -i[\mathbf{H}_{\Lambda}, \rho] + \sum_{\mu=1}^{d^k} 2L_{\Lambda,\mu}\rho L_{\Lambda,\mu}^{\dagger} - \{L_{\Lambda,\mu}^{\dagger}L_{\Lambda,\mu}, \rho\}$$

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Assumptions:

- N sites with finite local dimension d

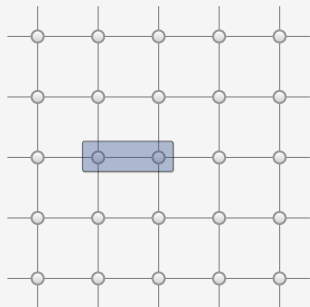


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Assumptions:

- N sites with finite local dimension d
- k -locality

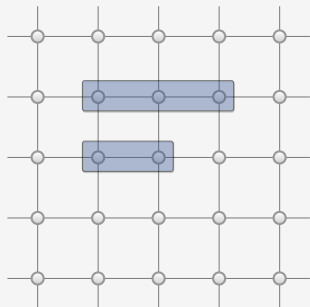


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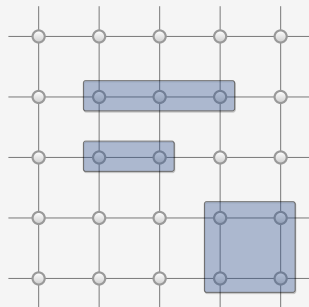


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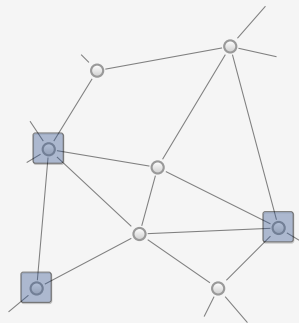


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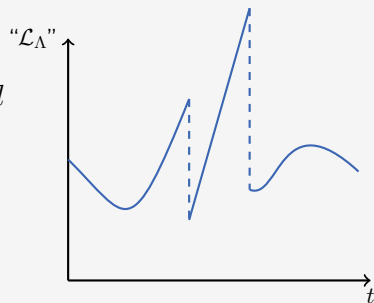


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Assumptions:

- N sites with finite local dimension d
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- arbitrary time dependence

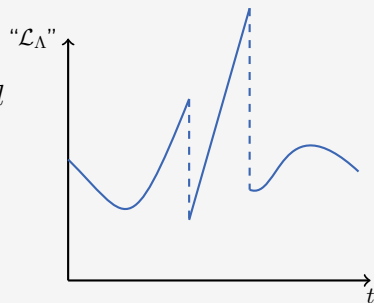


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Assumptions:

- N sites with finite local dimension d
- k -locality
- arbitrary time dependence
- $\|\mathbf{H}_{\Lambda}\|_{\infty}$ and $\|L_{\Lambda,\mu}\|_{\infty}$ bounded independent of N



Trotterization – what it is and how we get there

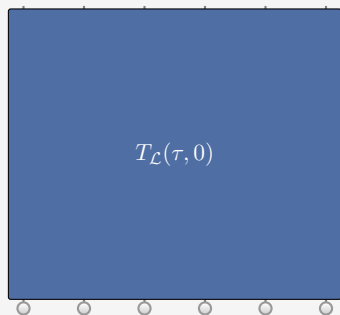
What we are aiming for:

$$T_{\mathcal{L}}(\tau, 0) \approx \prod_{j=1}^m \prod_{\Lambda}^K T_{\mathcal{L}_{\Lambda}}(\Delta t j, \Delta t(j-1))$$

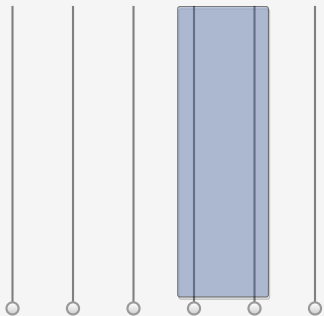
What we have to do:

- 1 Decompose $T_{\mathcal{L}}(\tau, 0)$ in time slices.
- 2 Approximate each time slice by applying local Liouvillians sequentially.

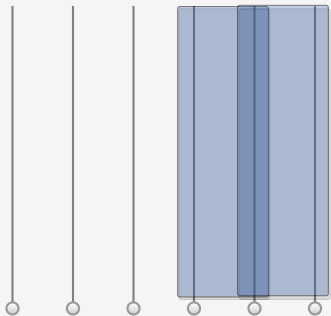
Trotterization of k -local Liouvillian dynamics



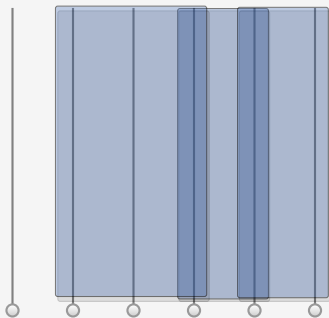
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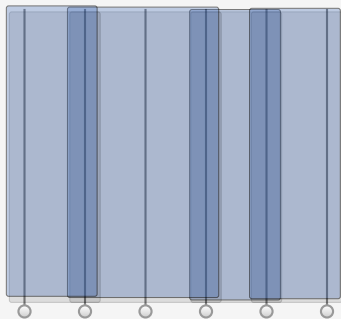
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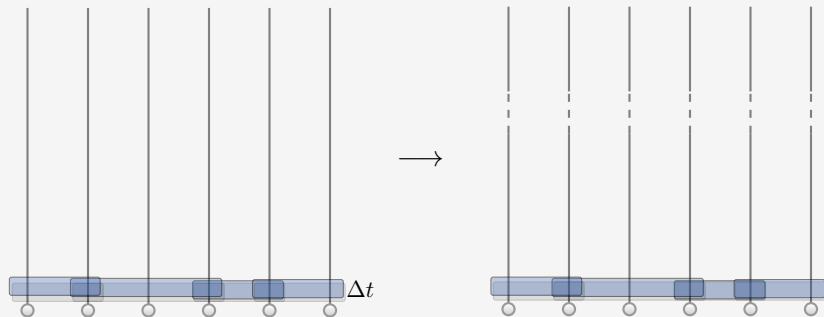
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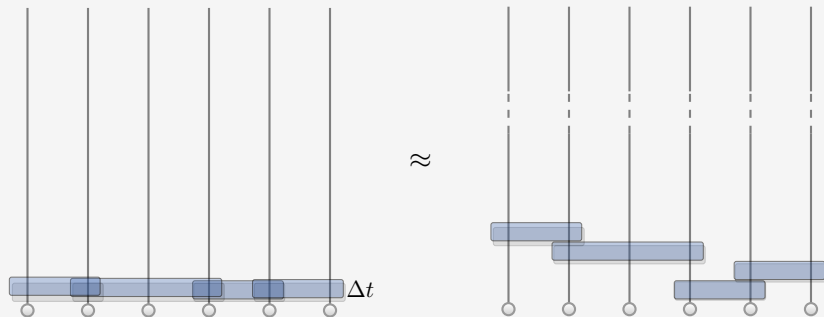
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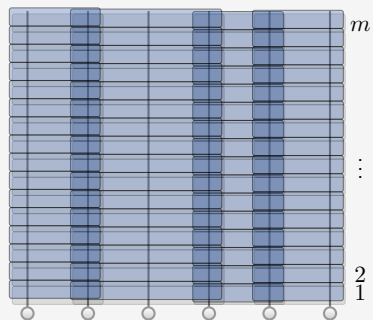
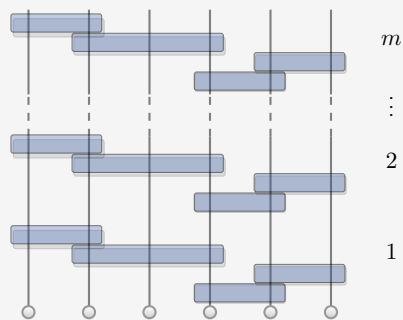
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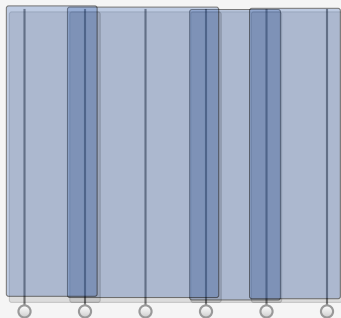
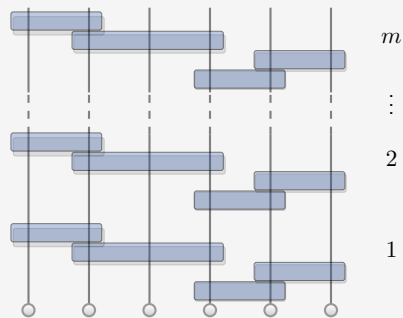
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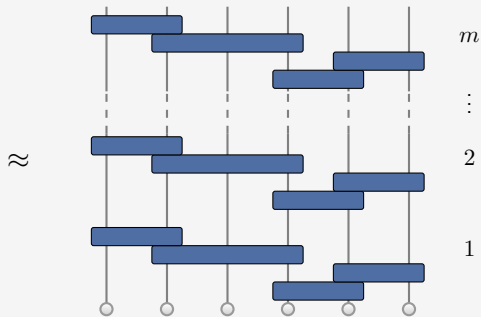
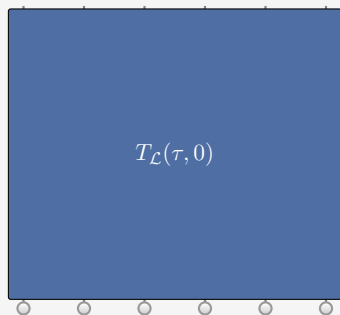
Trotterization of k -local Liouvillian dynamics

 \approx 

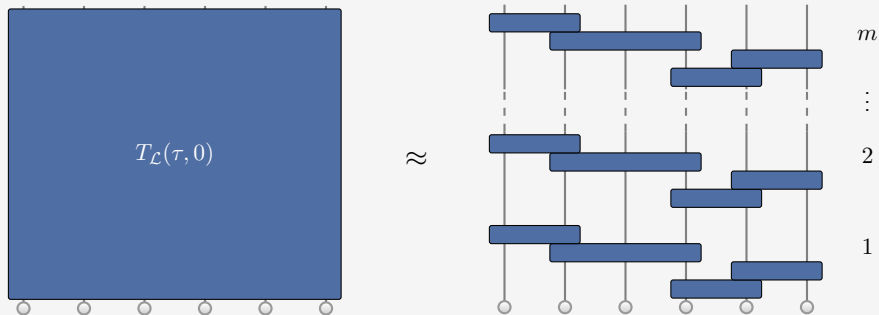
Trotterization of k -local Liouvillian dynamics


 \approx


Trotterization of k -local Liouvillian dynamics



Trotterization of k -local Liouvillian dynamics



$$m = \mathcal{O}\left(\frac{d^{2k} K^2 \tau^2}{\epsilon}\right)$$

A Trotter Formula for Liouvillian dynamics

Theorem 1 (Trotter decomposition [2])

Let $\mathcal{L} = \sum_{\Lambda}^K \mathcal{L}_{\Lambda}$ be a *piecewise continuous time dependent, k -local Liouvillian* acting on N subsystems of dimension d , then

$$\left\| T_{\mathcal{L}}(\tau, 0) - \prod_{j=1}^m \prod_{\Lambda} T_{\mathcal{L}_{\Lambda}}(\Delta t j, \Delta t(j-1)) \right\|_{1 \rightarrow 1} \leq c K^2 \tau \Delta t e^{b \Delta t},$$

with $b \in O(d^k)$, $c \in O(d^{2k})$.

A Trotter Formula for Liouvillian dynamics

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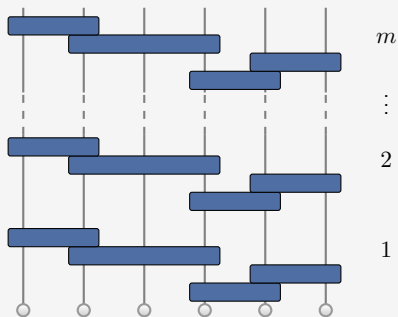
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$$\left\| T_{\mathcal{L}}(\tau, 0) - \prod_{j=1}^m \prod_{\Lambda} T_{\mathcal{L}_{\Lambda}^{\text{av}}}(\Delta t j, \Delta t (j-1)) \right\|_{1 \rightarrow 1} \leq c K^2 \tau \Delta t e^{b \Delta t},$$

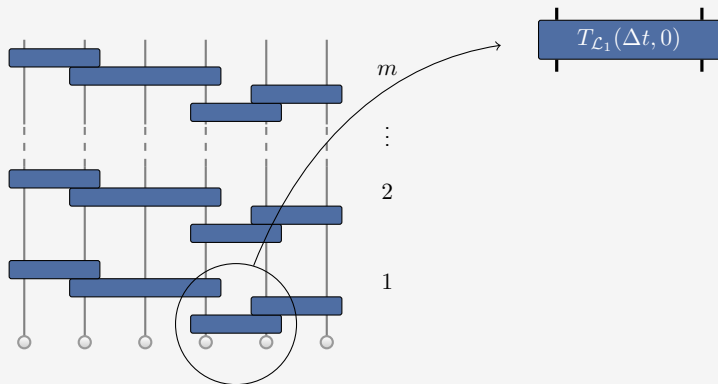
with $b \in O(d^k)$, $c \in O(d^{2k})$.

$$T_{\mathcal{L}_{\Lambda}^{\text{av}}}(\Delta t j, \Delta t (j-1)) = \exp(\Delta t \mathcal{L}_{\Lambda}^{\text{av}}) \quad \mathcal{L}_{\Lambda}^{\text{av}} = \Delta t \int_{\Delta t (j-1)}^{\Delta t j} \mathcal{L}_{\Lambda} dt$$

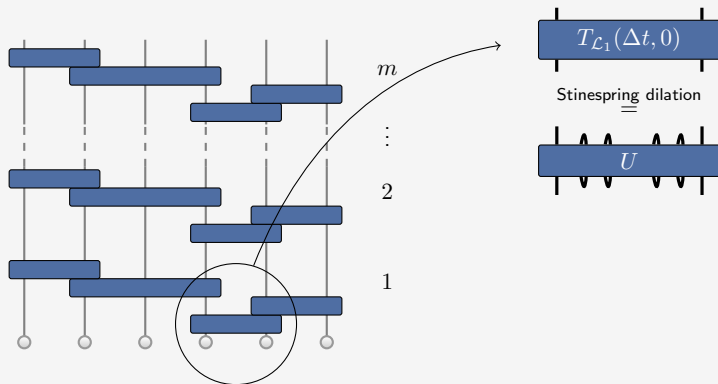
Stinespring and Solovay-Kitaev



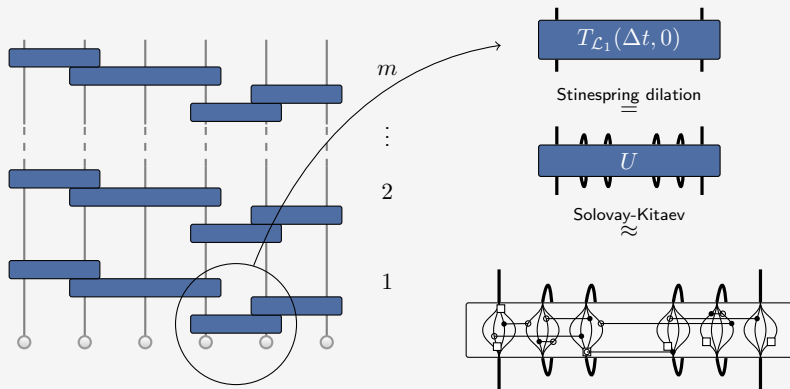
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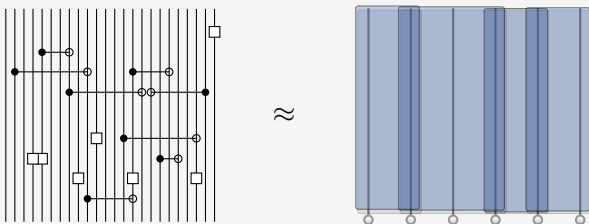


Implications

Implication 1

Power of dissipative quantum computing [3, 2]

Dissipative quantum computing with k -local, arbitrary time dependent Liouvillian dynamics is exactly as powerful as the circuit model.



[3] F. Verstraete, M. Wolf, and I. Cirac, Nature Physics, Vol 5, (2009) 633

[2] M. Kliesch, et al., PRL 107 (2011) 120501

Implication 2

Limits on efficient state preparation [2]

Even with arbitrary time-dependent k -local Liouvillian dynamics one can only reach exponentially few states after polynomial time.

[4] D. Poulin, A. Qarry, R. Somma, and F. Verstraete, PRL 106 (2011) 170501

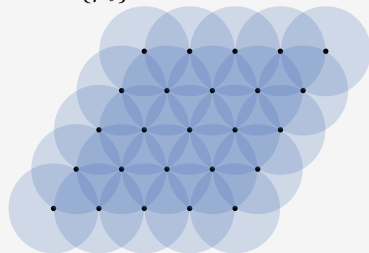
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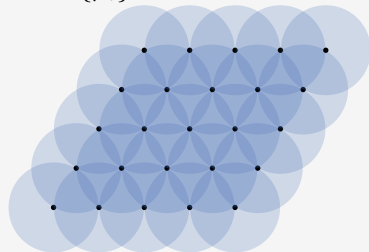
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$$\Omega(\exp(d^N))$$

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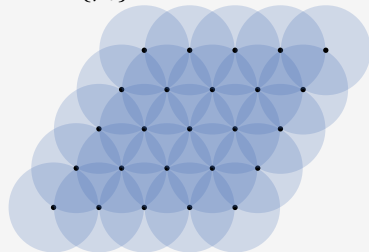
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Number of circuits for
 ϵ -approximation:

$$O\left(\exp(N^{3k+2}\tau^4)\right)$$

[4] D. Poulin, A. Qarry, R. Somma, and F. Verstraete, PRL 106 (2011) 170501

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Implication 3

Simulation on classical computers [2]

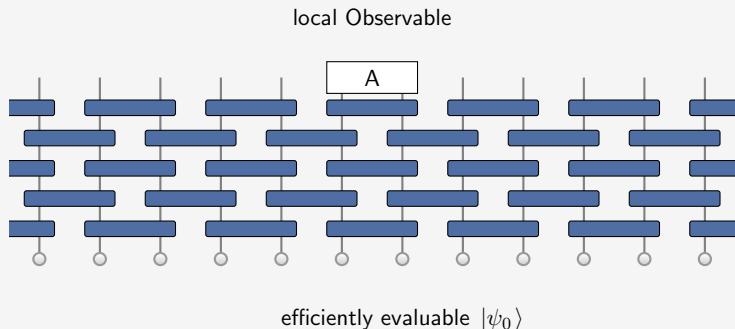
For fixed τ dissipative dynamics can be simulated efficiently in N on classical computers.

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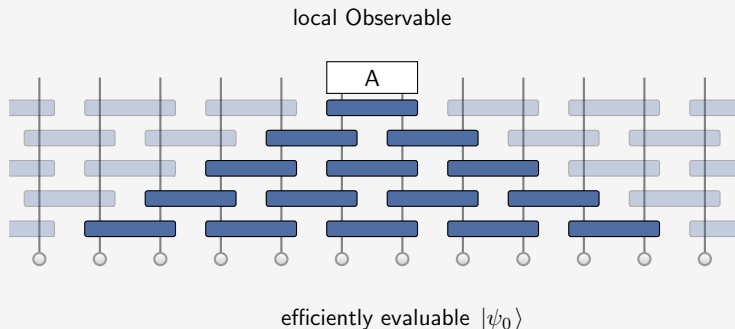


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For fixed τ dissipative dynamics can be **simulated efficiently in N** on classical computers.



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Implication 4

Strong quantum Church-Turing thesis [2]

Every quantum mechanical process that can be thought of as a computation can be efficiently simulated in the unitary circuit model of quantum computing.

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Implication 4

Strong quantum Church-Turing thesis [2]

Every quantum mechanical process that can be thought of as a computation can be efficiently simulated in the unitary circuit model of quantum computing.

Remember the **assumption** we made:

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Arguably the **most broadest setting** that allows **efficient simulation**.

[2] M. Kliesch, et al., PRL 107 (2011) 120501

Summary

$$T_{\mathcal{L}}(\tau, 0) \approx \prod_{j=1}^m \prod_{\Lambda}^K T_{\mathcal{L}_{\Lambda}}(\Delta t j, \Delta t (j - 1))$$

- k -local Liouvillian dynamics can be trotterized

Summary

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- Dissipative quantum computing is no more powerful than the circuit model
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- k -local Liouvillian dynamics can be simulated classically (efficient in N , inefficient in τ)

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- k -local Liouvillian dynamics can be trotterized
- Dissipative quantum computing is no more powerful than the circuit model
- Most states can not be prepared efficiently
- k -local Liouvillian dynamics can be simulated classically (efficient in N , inefficient in τ)
- A strong quantum Church-Turing theorem holds

Collaborators



Martin Kliesch



Thomas Barthel



Jens Eisert



Michael Kastoryano

References

Thank you for your attention!

→ slides: www.cgogolin.de

- [1] A. M. Turing,
"On computable numbers, with an application to the entscheidungsproblem",
Proc. London Math. Soc. **42** (1937) no. 230, 230–265.
- [2] M. Kliesch, T. Barthel, C. Gogolin, M. Kastoryano, and J. Eisert,
"Dissipative Quantum Church-Turing Theorem",
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- [3] F. Verstraete, M. M. Wolf, and J. Ignacio Cirac,
"Quantum computation and quantum-state engineering driven by dissipation",
Nature Physics **5** (2009) no. 9, 633.
- [4] D. Poulin, A. Qarry, R. Somma, and F. Verstraete,
"Quantum Simulation of Time-Dependent Hamiltonians and the Convenient Illusion of Hilbert Space",
Physical Review Letters **106** (2011) no. 17, 170501.